

Lecture five

Topics that must be covered in this lecture:

- *Finite state automata with empty move.*
- *Equivalence of NFA with and without ϵ move*

Finite state automata with empty move ϵ :

If the definition of a NFA is altered, so that moves from one state to another can be accomplished without any input we say that the automaton has ϵ -moves. NFA accepts W if there is some path labeled W from initial state to a final state. Of course, edges labeled ϵ may be included in the path, although the ϵ 's do not appear explicitly in W .

More formally, a NFA $M=(Q, \Sigma, \delta, S, F)$ has ϵ -moves if δ , instead of being a function $Q \times \Sigma \rightarrow 2^Q$, is defined as a function $Q \times (\Sigma \cup \{\epsilon\}) \rightarrow 2^Q$.

$$\delta'(Q, \epsilon) = R(Q) \text{ or } \epsilon\text{-closure}(Q)$$

$$\delta'(Q, \Sigma) = R(\delta(R(Q), \Sigma))$$

Equivalence of NFA with and without ϵ move:

If there is NFA with ϵ , $M = (Q, \Sigma, \delta, q_0, F)$ then there is NFA without empty move $M' = (Q', \Sigma, \delta', q_0', F')$

Define $\delta' : Q \times (\Sigma \cup \{\epsilon\}) \rightarrow Q$

By:

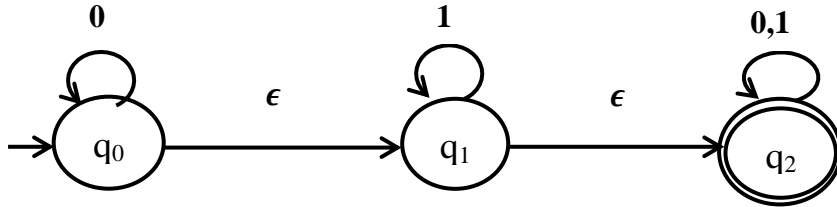
$$\delta'(K, \epsilon) = R(K) \text{ or } \epsilon\text{-closure}(K)$$

$$\delta'(K, a) = R(\delta(R(K), a))$$

The set of final states F' is $F \cup \{q\}$ if $R(q) \in F$

Example1:

Convert the NFA with empty move into NFA without empty move.



Solution:

Note: 1- ϵ - closure (ϵ^*): All the states that can be reached from a particular state only by seeing the ϵ symbol.

2- Every state on ϵ goes to itself.

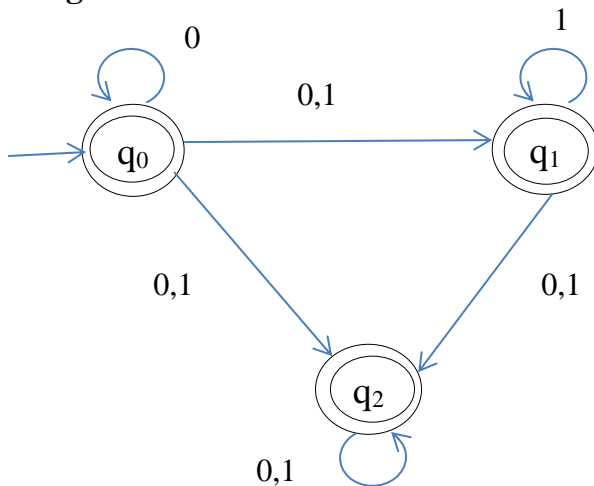
	ϵ^*	0	ϵ^*
q0	q0 q1 q2	q0 Φ q2	q0,q1, q2 - q2
q1	q1 q2	Φ q2	- q2
q2	q2	q2	q2

	ϵ^*	1	ϵ^*
q0	q0 q1 q2	Φ q1 q2	- q1,q2 q2
q1	q1 q2	q1 q2	q1,q2 q2
q2	q2	q2	q2

Transition table of NFA without ϵ :

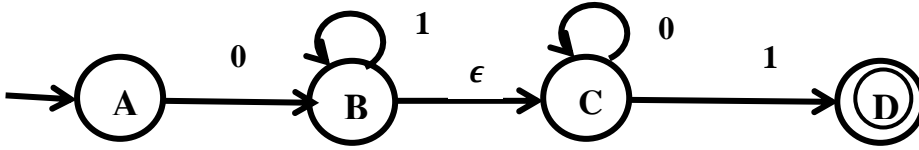
	0	1
q0	{q0,q1,q2}	{q1,q2}
q1	{q2}	{q1,q2}
q2	{q2}	{q2}

The state diagram of NFA without ϵ :



Example2:

Convert the NFA with empty move into NFA without empty move:



Solution:

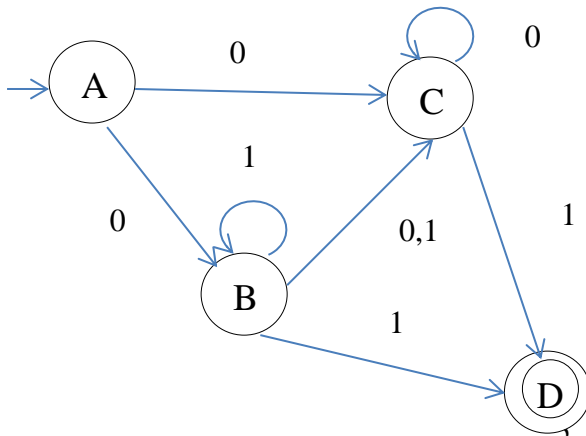
	ϵ^*	0	ϵ^*
A	A	B	B,C
B	B C	Φ C	- C
C	C	C	C
D	D	Φ	-

	ϵ^*	1	ϵ^*
A	A	Φ	-
B	B C	B D	B,C D
C	C	D	D
D	D	Φ	-

Transition table of NFA without ϵ :

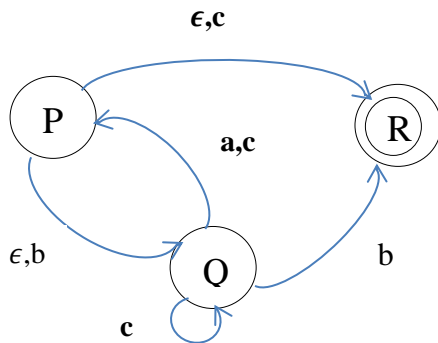
	0	1
A	{B,C}	Φ
B	{C}	{B,C,D}
C	{C}	{D}
D	Φ	Φ

The state diagram of NFA without ϵ :



Example3:

Convert the NFA with empty move into NFA without empty move:



Solution:

	ϵ^*	a	ϵ^*
P	P Q R	Φ P Q	- P,Q,R -
Q	Q	P	P,Q,R
R	R	Φ	-

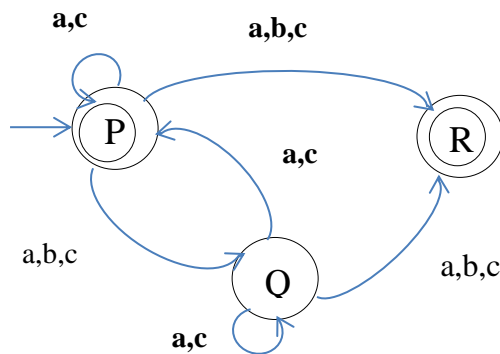
	ϵ^*	b	ϵ^*
P	P Q R	Q R Φ	Q R -
Q	Q	R	R
R	R	Q	-

	ϵ^*	c	ϵ^*
P	P Q R	R Q,P Φ	R P,Q,R -
Q	Q	Q	P,Q,R
R	R	Φ	-

Transition table of NFA without ϵ :

	a	b	c
P	{P,R,Q}	{Q,R}	{P,Q,R}
Q	{P,R,Q}	{R}	{P,Q,R}
R	Φ	Φ	Φ

The state diagram of NFA without ϵ :

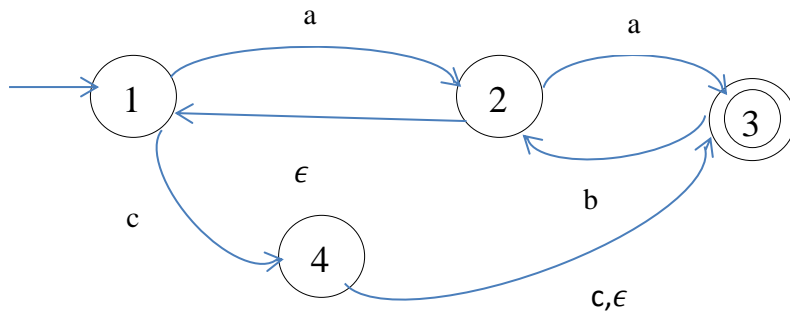


Conversion NFA with ϵ to DFA:

To construct DFA equivalent to the NFA with empty move apply the following steps:

- 1- Draw the NFA transition table and add new column contains ϵ -closure for all states in the NFA.
- 2- Start computing the DFA transition table from the first state and take the resulting states as the next state in each step.

Example1: construct DFA equivalent to the following NFA with empty move:



Solution:

	ϵ^*	a	ϵ^*
1	1	2	1,2
2	1 2	2 Φ	1,2 -
3	3	2	1,2
4	3 4	2 Φ	1,2 -

	ϵ^*	b	ϵ^*
1	1	Φ	-
2	1 2	Φ 3	- 3
3	3	Φ	-
4	3 4	Φ Φ	- -

	ϵ^*	c	ϵ^*
1	1	4	3,4
2	1 2	4 Φ	3,4
3	3	Φ	-
4	3 4	Φ 3	- 3

1- The transition table of NFA:

	a	b	c
1	1,2	Φ	3,4
2	1,2	3	3,4
3	1,2	Φ	Φ
4	1,2	Φ	3

2- The transition table of DFA:

	a	b	c
-1	1,2	D	3,4
1,2	1,2	3	3,4
+3,4	1,2	D	3
+3	1,2	D	D
D	D	D	D

3- The state diagram for DFA:

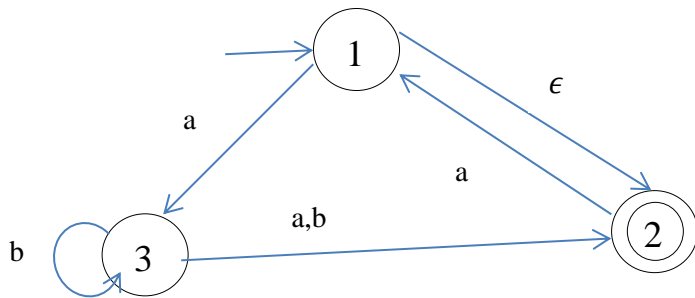
طريقة اخرى:**Solution:****1- The transition table of NFA with empty move :**

	a	b	c	ϵ -closure
1	2	-	4	1
2	-	3	-	2,1
3	2	-	-	3
4	-	-	3	4,3

2- The transition table of DFA:

	a ϵ^*	b ϵ^*	c ϵ^*
-1	2,1	D	4,3
2,1	2,1	3	4,3
+4,3	2,1	D	3
+3	2,1	D	D
D	D	D	D

Example2: construct DFA equivalent to the following NFA with empty move:



Solution:

1- The transition table of NFA:

	a	b	ϵ -closure
1	3	-	1,2
2	1	-	2
3	2	2,3	3

2- The transition table of DFA:

	a ϵ^*	b ϵ^*
-,+ 1	3	D
3	2	2,3
+2	1,2	D
+2,3	1,2	2,3
+1,2	1,2,3	D
+1,2,3	1,2,3	2,3
D	D	D

3-The state diagram for DFA: