

TIKRIT UNIVERSITY
COLLEGE OF COMPUTER SCIENCE AND MATHEMATICS
DEPARTMENT OF COMPUTER SCIENCE



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Lexical Errors

Type of errors for Lexical analysis

Lexical errors include misspellings of identifiers, keywords, or operators... Etc.

Ex:

fi (a == f(x))

Problem of lexical analyzer cannot tell:

- Whether **fi** is a misspelling of the keyword **if** .
- Or **fi** is an undeclared function identifier .

Terms Tokens, Patterns, and Lexemes

Source program

Ex.:
`if (x > y)`
`x = 10 ;`
`while ...`

Lexical analyzer phase

1

Lexemes
if
(
x
>
y
)
x
=
10
;
while

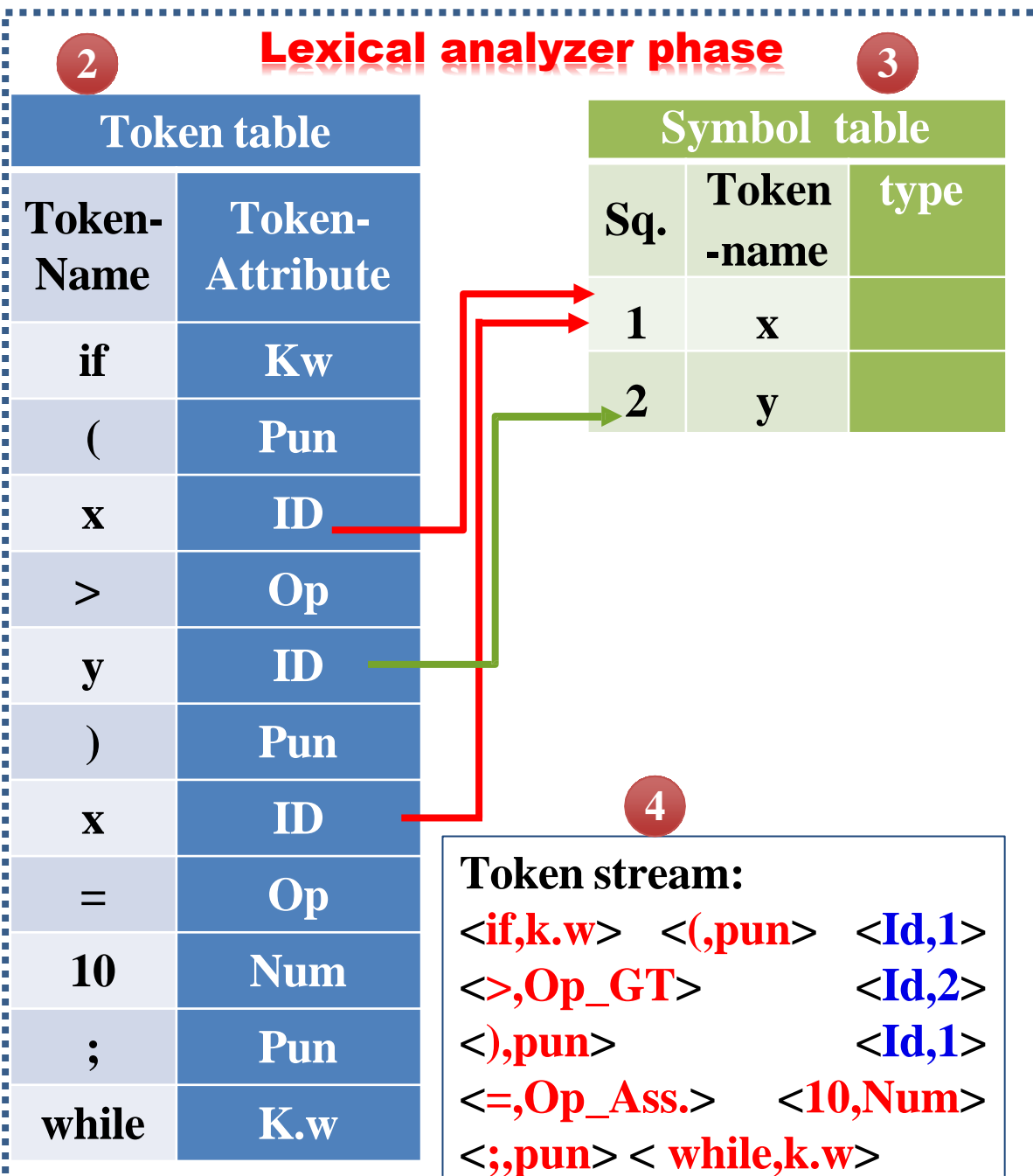
(1) Lexeme:

A lexeme: is one character or a sequence of characters in the source program that denotes the meaning a basic lexical unit of the source program and is identified (breaks up) by the lexical analyzer.

Source program

Ex.:

```
if ( x > y )
    x = 10 ;
while ...
```



(2) Token:

A token: is a pair consisting of a token name and an optional attribute value (token-name; attribute-value). The first component token-name is an abstract symbol that is used during syntax analysis, and the second component points to a kind of lexical unit, e.g., a particular keyword, or a sequence of input characters denoting an identifier.

Ex.:

```
if ( x > y )
    x = 10 ;
while ...
```

(3) Patterns:

Patterns: Each token has a pattern that describes which sequences of characters can form the lexemes corresponding to that token. The set of words, or strings of characters, that match a given pattern is called a language.

Pattern: is a description of a grammar of language that the tokens may take.

```
<Statement> → if ( <Expression> ) <Statement>
              | if ( <Expression> ) <Statement> else <Statement>
              | ε
<Expression> → <Term> relop <Term> | <Term>
<Term>       → id | number
<relop>      → < | > | <= | >= | == | !=
<Id>         → L ( L | D )*
<L>          → A | B | ... | Z | a | b | ... | z | _
<D>          → 0 | ... | 9
Number       → digits OptionalFraction OptionalExponent
<Ds>        → DD*
OptionalFraction → . Ds | ε
OptionalExponent → ( E ( + | - | ε Ds ) | ε
```

ملاحظة: القاعدة النحوية لتعريف **if** سوف يولد مشكلة التداخل من اليسار لذلك يجب ان تكتب بالصيغة التالية:

$stmt \rightarrow if (expr) stmt [else stmt]? / \epsilon$

Specification of Tokens

Alphabet, Strings and Languages

1. Alphabet: is any finite set of symbols.

- Binary alphabet $\longrightarrow \{0,1\}$
- Digit alphabet $\longrightarrow \{0,1,2,3,4,5,6,7,8,9\}$
- Letter alphabet $\longrightarrow \{a,b,c,\dots,z, A,B,C,\dots,Z\}$
- ASCII alphabet: a set of digital codes (0 , 1) is used in many software systems.
- Unicode alphabet include 100000 characters.

ASCII: American Standard Code for Information Interchange, a set of digital codes (0 , 1) representing letters, numerals, and other symbols, widely used as a standard format in the transfer of text between computers.

Unicode: an international encoding standard for use with different languages and scripts, by which each letter, digit, or symbol is assigned a unique numeric value.

Alphabet, Strings and Languages

2. String ("Sentence" and "Word"): A string over an alphabet is a finite sequence of symbols drawn from the alphabet.

- The length of a string **s**, usually written **|s|**, is the number of occurrences of symbols in **s**.

010111011100 is string from the binary alphabet. The string of length **12**.

01 is string from the binary alphabet. The string of length **2**.

0 is string from the binary alphabet. The string of length **1**.

banana is string from the Letter alphabet. The string of length **6**.

ϵ the empty string , the string of length **zero**.

Specification of Tokens

Alphabet, Strings and Languages

3. Language: is any countable set of strings over some fixed alphabet.

Strings	Language
010111011100 01 0	Ex1: This is machine Language.
{ if (x==y) x = x + 1; for (int i =0 ; i<10; i++) y = y + I; }	Ex2: This is source program of C++ Language.
<{,Pun> <if,k.w> <(, Pun > <Id,1 pointer to symbol-table entry for x> <==,Op> <Id,2 pointer to symbol-table entry for y) <(, Pun > ... < , >	Ex3: This is Language of token stream by C++.
The empty set, or $\{\epsilon\}$, the set containing only the empty string.	Abstract languages like \emptyset .
The set of all grammatically correct English sentences.	English Language.

Terms for Parts of Strings

(9)

1. A **prefix** of string **s** is any string obtained by removing zero or more symbols from the end of **s**.

Ex: **ban**, **banana**, and ϵ are prefixes of **banana**.

✓ اي جزء من السلسلة من البداية.
✓ السلسلة نفسها.
✓ ϵ .

2. A **suffix** of string **s** is any string obtained by removing zero or more symbols from the beginning of **s**.

Ex: **nana**, **banana**, and ϵ are suffixes of **banana**.

✓ اي جزء من السلسلة من النهاية.
✓ السلسلة نفسها.
✓ ϵ .

3. A **substring** of **s** is obtained by deleting any prefix and any suffix from **s**.

Ex: **banana**, **anan**, **nan**, and ϵ are substrings of **banana**.

✓ اي جزء من السلسلة من أي مكان بعد حذف رمز البداية والنهاية.
✓ والسلسلة نفسها.
✓ ϵ .

Terms for Parts of Strings

(10)

4. The **proper prefixes, suffixes, and substrings** of a string **s** are those, prefixes, suffixes, and substrings, respectively, of **s** that are not ϵ or not equal to **s** itself.

Ex: **ba, ban,** and **banan** are proper prefixes of **banana**.

Ex: **na, nana** and **anana** are proper suffixes of **banana**.

Ex: **anan, nan** are proper substrings of **banana**.

✓ اي جزء من السلسلة من أي مكان بشرط ان تكون رموز متعاقبة.

✓ لا يشمل السلسلة نفسها.

✓ لا يشمل ϵ .

5. A **subsequence** of **s** is any string formed by deleting zero or more not necessary consecutive positions of **s**.

Ex: **baan** is a subsequence of **banana**.

✓ اي جزء من السلسلة ليس من الضروري متعاقب.

✓ يشمل السلسلة نفسها.

✓ يشمل ϵ .

Concatenation: التتابع

- If **x** and **y** are strings,
- Then the concatenation of **x** and **y**, denoted **xy**,
- is the string formed by appending **y** to **x**.
- Ex: if **x = dog** and **y = house**, then **xy = doghouse**.
- The empty string is concatenation.
- Any string **s**, ϵs
 $= s\epsilon = s$.

Exponentiation : التسلسل

Exponentiation of strings is a product of concatenation.

Since:

$$s^0 = \varepsilon$$

$$\varepsilon s = s \varepsilon = s$$

define: s^0 to be ε

$$s^1 = s$$

$$s^2 = ss$$

$$s^3 = sss$$

$i > 0$, define s^i to be $s^{i-1}s$.

الأُس في السلاسل
هي
انتاج من تسلسل

OPERATION	DEFINITION AND NOTATION
<i>Union of L and M</i>	$L \cup M = \{s \mid s \text{ is in } L \text{ or } s \text{ is in } M\}$
<i>Concatenation of L and M</i>	$LM = \{st \mid s \text{ is in } L \text{ and } t \text{ is in } M\}$
<i>Kleene closure of L</i>	$L^* = \bigcup_{i=0}^{\infty} L^i$
<i>Positive closure of L</i>	$L^+ = \bigcup_{i=1}^{\infty} L^i$

Figure 3.6: Definitions of operations on languages

1. **LUD** is the set of letters and digits. **62** strings of length one.

$$\begin{array}{l} \mathbf{A | B | \dots | Z | a | b | \dots | z} = \mathbf{52} \\ \mathbf{0 | 1 | \dots | 9} = \mathbf{10} \end{array} \quad \left. \vphantom{\begin{array}{l} \mathbf{A | B | \dots | Z | a | b | \dots | z} \\ \mathbf{0 | 1 | \dots | 9} \end{array}} \right\} \mathbf{62}$$

2. **LD** is the set of **520** strings of length two.

A0	B0	...	Z0	a0	b0	...	z0
A1	B1	...	Z1	a1	b1	...	z1
A2	B2	...	Z2	a2	b2	...	z2
...
...
...
A9	B9	...	Z9	a9	b9	...	z9
Total 520							

3. L^4 is the set of all 4 -letter strings.
4. L^* is the set of all strings of letters, including ϵ the empty string.
5. $L(LUD)^*$ is the set of all strings of letters and digits beginning with a letter.
6. D^+ is the set of all strings of one or more digits.
7. ? Zero or one instance.

THANK YOU